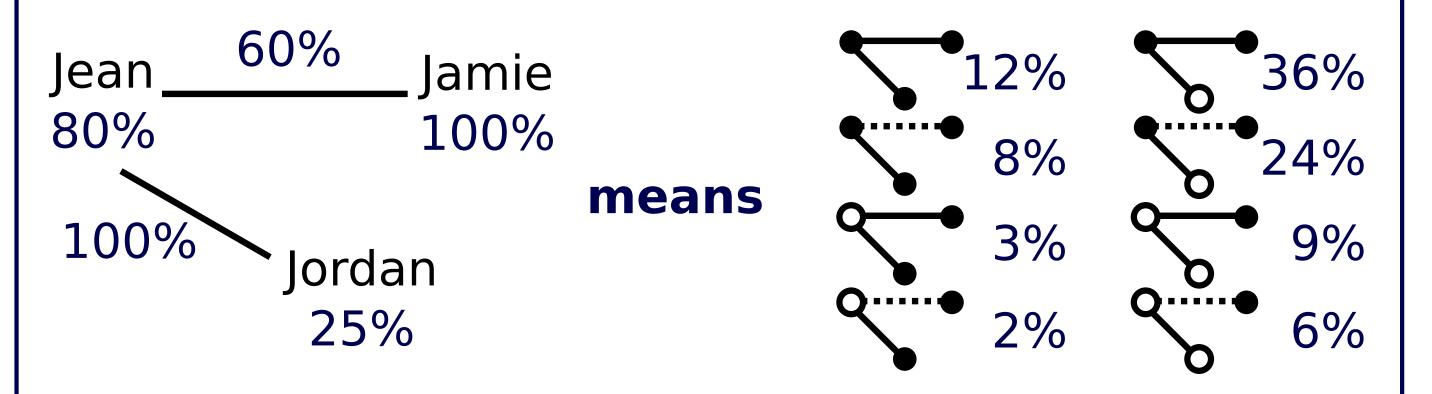
Leveraging the Structure of Uncertain Data Antoine Amarilli Télécom ParisTech, CNRS LTCI, Université Paris-Saclay joint work with Pierre Bourhis and Pierre Senellart

Probabilistic Databases

A dating website uses machine learning to classify users as active/inactive, and chats as flirtatious or not

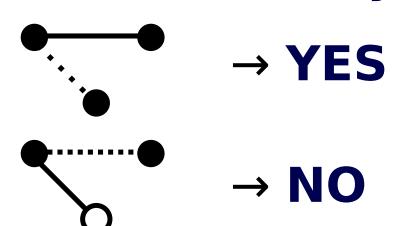


We assume **independence** across all these facts

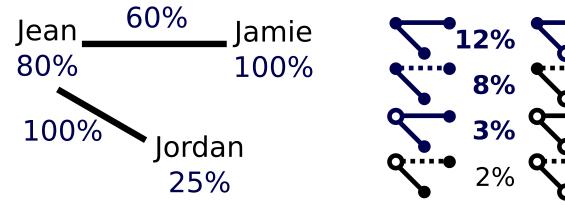
Query Evaluation

Query: are there **active users** engaged in **flirtatious** chat?

On **deterministic** data, this is **easy**:



On **probabilistic** data, evaluation becomes much harder!





Can we compute **efficiently** the probability of a **query**?

→ The task is **intractable** (#P-hard) for many queries

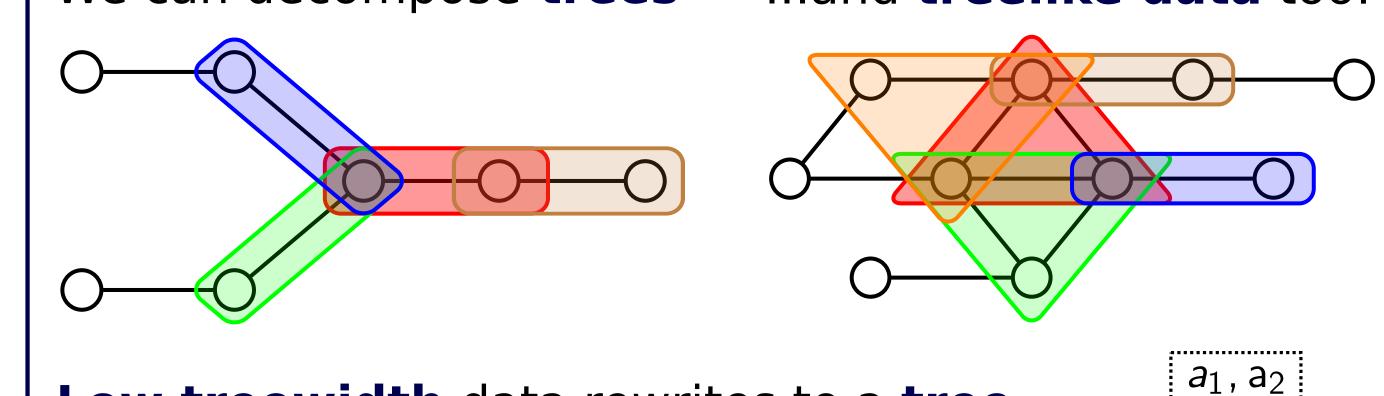
 \rightarrow Cannot represent **correlations**, like "Monoamorous users are flirtatious with ≤ 1 person"

even when the query is fixed (i.e., in **data complexity**) → What can we do? (especially on simple, realistic data?)

Problem Statement: Which structural hypotheses on probabilistic data make query evaluation tractable?

Treewidth

A **measure** of how much the data is similar to a **tree** We can decompose **trees** ...and **treelike data** too:



Low treewidth data rewrites to a tree

Lineages

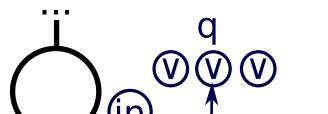
Lineage φ of an arbitrary Boolean **query** q on **database** D: φ is a **Boolean formula** (or **circuit**) on the facts of D such that $D' \subseteq D$ satisfies q iff φ holds for the valuation of D'

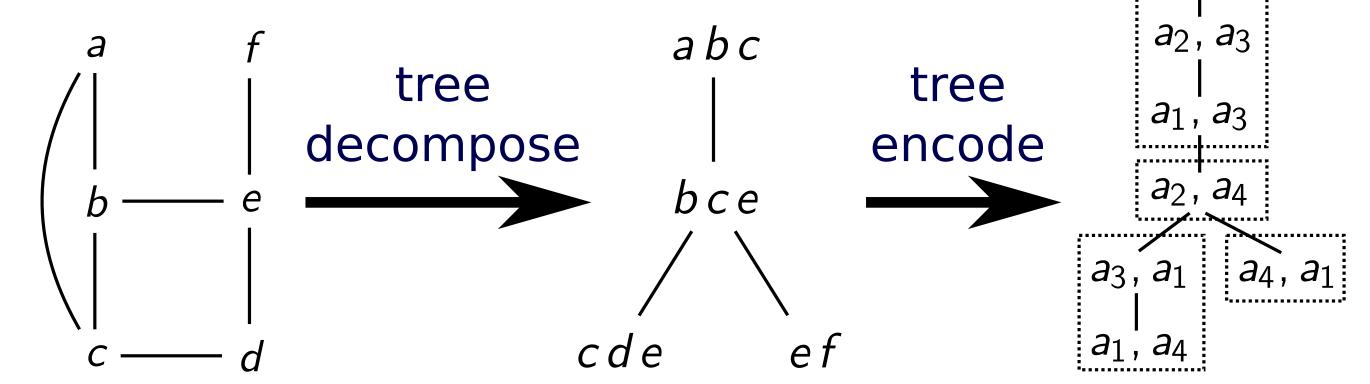
Jean $\frac{m_{ea}}{j_e}$ Jamie j_a m_{eo} Jordan

Is there a flirtatious pair of active users? **Lineage:** $j_e \land ((m_{ea} \land j_a) \lor (m_{eo} \land j_o))$

Our main technical results:

so probability computation is 1





Courcelle's theorem: Monadic second-order queries are tractable to evaluate on non-probabilistic data using tree automata on the tree encoding

Results

Our **main result**:

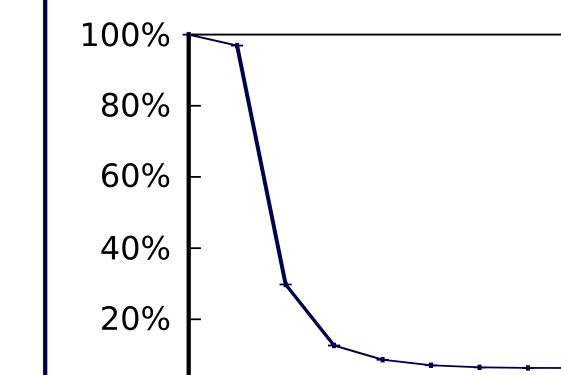
For any monadic second-order **query** q and integer k for any input TID **database** D of treewidth $\leq k$ we can compute in O(|D|) the **probability** of q on D (up to polynomial arithmetic costs)

- Lineage circuits for tree automata on (in)... **uncertain trees** can be computed in O(n) - Extends to **bounded treewidth** data \mathbf{q}_2 \mathbf{q}_1 OOC 000 - The circuit has **low treewidth** itself (in) (in)

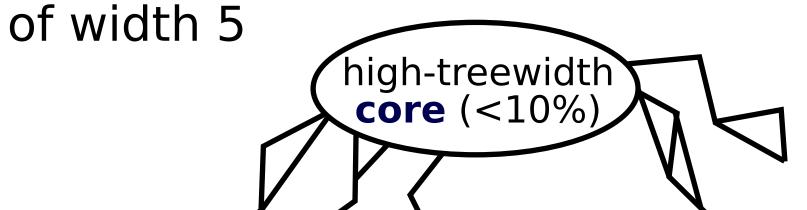
Ongoing Work with Mikaël Monet, Silviu Maniu, and Pierre Senellart

• Computing tree decompositions on real datasets

The Paris **road network** (4.3M nodes and 5.4M edges) has treewidth ≤ 521 (computed using heuristics)



Most of the network is covered by a partial decomposition



Extensions:

 \rightarrow Also for low-treewidth **correlations** (e.g. **mutually exclusive** database facts) \rightarrow Also for expressive **lineages** ($\mathbb{N}[X]$ -provenance) **Lower bound:** Query evaluation on TID is **intractable**

if the treewidth is not bounded (under some technical conditions)

References

Antoine Amarilli, Pierre Bourhis, Pierre Senellart: - Provenance Circuits for Trees and Treelike Instances, ICALP'15 - Tractable Lineages on Treelike Instances: Limits and Extensions, PODS'16

low-treewidth **tentacles** (>90%) 0% width 10 width 5 \rightarrow Answer queries with **uncertainty** (e.g., RER trip time) with tree automata on the tentacles and sampling on the core

- Tractability in **combined complexity** for restricted queries Our method on low-treewidth data is **linear** in the data but hides high complexity in the **query**
- \rightarrow Lower bound: Tractable complexity in TID data and query seems unlikely, even for tree-shaped queries and data

 \rightarrow Can we compute **lineages** more efficiently in some cases?