

# Smoothing Structured Decomposable Circuits

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# Probabilistic Circuits

Tractable computation graph, encoding a distribution.

SOTA for:

- ▶ Inference algorithms for PGMs
- ▶ Inference algorithms for probabilistic programs
- ▶ Discrete density estimation

Check out:

**Tractable Probabilistic Models:** (UAI19 / AAAI20 tutorial)

# Tractability

Different combination of properties leads to different families of circuits

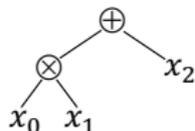
	SPN	AC	PSDD
Decomposability	✓	✓	✓(S)
Determinism	X	✓	✓
Smoothness	✓	✓	✓
Pr(evid)	✓	✓	✓
Marginal	✓	✓	✓
MPE	X	✓	✓
Marginal MAP	X	X	✓*
Expectation	X	X	✓*

...with different tractability properties.

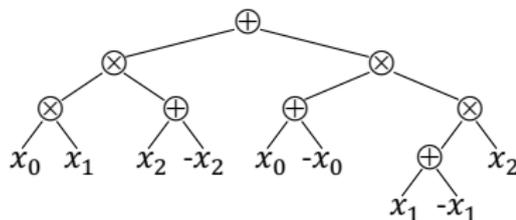
# Smoothness

## Definition

A circuit is **smooth** if for every pair of children  $c_1$  and  $c_2$  of a  $\oplus$ -gate,  $\text{vars}_{c_1} = \text{vars}_{c_2}$ .



(a) A circuit.



(b) A smooth circuit.

**Figure:** Two equivalent circuits computing  $(x_0 \otimes x_1) \oplus x_2$ . The left one is not smooth and the right one is smooth.

# Smoothing a Circuit: Naive Quadratic Algorithm

- ▶ Go to each gate  $O(m)$  and fill in each variable  $O(n)$
- ▶ Complexity  $O(nm)$
- ▶ Problematic when  $n \geq 1,000$  and  $m \geq 1,000,000$

Our near-linear smoothing algorithm:  $O(m \cdot \alpha(m, n))$

# Smoothing a Circuit: Missing Intervals

Key Insight: missing variables for each gate form two intervals (in the inorder traversal of vtree).

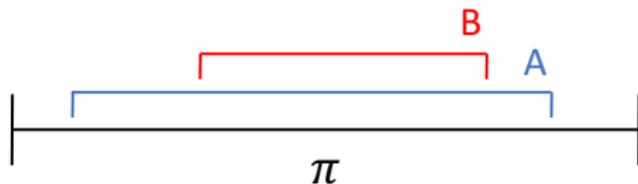


Figure:  $A \setminus B$  forms two intervals

We need to fill in  $2m$  intervals.

# Semigroup Range Sum

## Theorem

*Given  $n$  variables defined over a semigroup and  $m$  intervals, the sum of all intervals can be computed using  $O(m \cdot \alpha(m, n))$  additions [Chazelle and Rosenberg 1989].*

$\alpha(m, n)$  is the inverse Ackermann function, which grows very slowly.

\*The original theorem only bounds the number of additions. We bound the number of total operations.

# Takeaways

- ▶ Probabilistic circuits can encode complex distributions.
- ▶ They can compute exact likelihoods, marginals, and more
  - ▶ But **only if they are smooth**.
- ▶ Best smoothing algorithm was quadratic.
- ▶ We propose a near **linear time** smoothing algorithm.

# Thanks!

Poster: East Exhibition Hall B+C #182, 10:45AM

Code: <https://github.com/AndyShih12/SSDC>

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